Aspects of Tree-Based Statistical Machine Translation

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Human Language Technology FBK-irst

2011

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Outline

Tree-based translation models:

- Synchronous context free grammars
- BTG alignment model
- Hierarchical phrase-based model

Decoding with SCFGs:

- Translation as Parsing
- DP-based chart decoding
- Integration of language model scores

Learning SCFGs

Rule extraction from phrase-tables

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Tree-Based Translation Models

Levels of Representation in Machine Translation:



- ▶ $\pi \mapsto \sigma$: tree-to-string
- $\sigma \mapsto \pi$: string-to-tree
- $\pi \mapsto \pi$: tree-to-tree

? Appropriate Levels of Representation ?

Tree Structures



Syntactic Structures:

- rooted ordered trees
- internal nodes labeled with syntactic categories
- leaf nodes labeled with words
- linear and hierarchical relations between nodes

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Tree-to-Tree Translation Models



- syntactic generalizations over pairs of languages: isomorphic trees
- syntactically informed unbounded reordering
- formalized as derivations in synchronous grammars

? Adequacy of Isomorphism Assumption ?

Context-Free Grammars

CFG (Chomsky, 1956):

- formal model of languages
- more expressive than FSAs and REs
- first used in linguistics to describe embedded and recursive structures

CFG Rules:

- left-hand side nonterminal symbol
- right-hand side string of nonterminal or terminal symbols

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 distinguished start nonterminal symbol

$$egin{array}{ccc} S o 0S1 & S ext{ rewrites as } 0S1 \ S o \epsilon & S ext{ rewrites as } \epsilon \end{array}$$

CFG Derivations

Generative Process:

- 1. Write down the **start** nonterminal symbol.
- 2. **Choose** a **rule** whose left-hand side is the left-most written down nonterminal symbol.
- 3. **Replace** the symbol with the right-hand side of that rule.
- 4. Repeat step 2 while there are nonterminal symbols written down.



CFG Formal Definitions

CFG $G = \langle V, \Sigma, R, S \rangle$:

- V: finite set of nonterminal symbols
- Σ: finite set of terminals, disjoint from V
- ▶ *R*: finite set of rule $\alpha \rightarrow \beta$, with α a nonterminal and β a string of terminals and nonterminals

► S: the start nonterminal symbol

Let u, v be strings of $V \cup \Sigma$, and $\alpha \rightarrow \beta \in R$, then we say:

- . $u\alpha v$ yields $u\beta v$, written as $u\alpha v \Rightarrow u\beta v$
- . *u* derives *v*, written as $u \Rightarrow^* v$, if u = v or a sequence u_1, u_2, \ldots, u_k exists for $k \ge 0$ and $u \Rightarrow u_1 \Rightarrow u_2 \Rightarrow \ldots \Rightarrow u_k \Rightarrow v$
- $L(G) = \{ w \in \Sigma^* | S \Rightarrow^* w \}$

CFG Examples

 G_1 :

$$R = \{S \rightarrow NP \ VP,$$

$$NP \rightarrow N | DET \ N | N \ PP,$$

$$VP \rightarrow V \ NP | V \ NP \ PP,$$

$$PP \rightarrow P \ NP,$$

$$DET \rightarrow the | a,$$

$$N \rightarrow Alice | Bob | trumpet,$$

$$V \rightarrow chased,$$

$$P \rightarrow with \}$$

G3:

$$\begin{split} R &= \{ NP \rightarrow NP \ CONJ \ NP | NP \ PP | DET \ N, \\ PP \rightarrow P \ NP, P \rightarrow of, \\ DET \rightarrow the | two | three, \\ N \rightarrow mother | pianists | singers, \\ CONJ \rightarrow and \} \end{split}$$

? **derivations** of the mother of three pianists and two singers

? derivations of

Alice chased Bob with the trumpet

- ► same parse tree can be derived in different ways (≠ order of rules)
- same sentence can have different parse trees (≠ choice of rules)

Transduction Grammars aka Synchronous Grammars

TG (Lewis and Stearns, 1968; Aho and Ullman, 1969):

- two or more strings derived simultaneously
- more powerful than FSTs
- used in NLP to model alignments, unbounded reordering, and mappings from surface forms to logical forms

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Synchronous Rules:

- left-hand side nonterminal symbol associated with source and target right-hand sides
- bijection [] mapping nonterminals in source and target of right-hand sides

$$\begin{cases} E \to E_{[1]} + E_{[3]} / + E_{[1]} E_{[3]} & \text{infix to Polish notation} \\ E \to E_{[1]} * E_{[2]} / * E_{[1]} E_{[2]} \\ E \to n / n & n \in N \end{cases}$$

Synchronous CFG

$$\begin{split} & \mathrm{NP} \longrightarrow \mathrm{DT}_{\overline{\mathrm{U}}} \mathrm{NPB}_{\overline{\mathrm{D}}} \; / \; \mathrm{DT}_{\overline{\mathrm{U}}} \mathrm{NPB}_{\overline{\mathrm{D}}} \\ & \mathrm{NPB} \longrightarrow \mathrm{JJ}_{\overline{\mathrm{U}}} \mathrm{NN}_{\overline{\mathrm{D}}} \; / \; \mathrm{JJ}_{\overline{\mathrm{U}}} \mathrm{NN}_{\overline{\mathrm{D}}} \\ & \mathrm{NPB} \longrightarrow \mathrm{NPB}_{\overline{\mathrm{U}}} \mathrm{JJ}_{\overline{\mathrm{D}}} \; / \; \mathrm{JJ}_{\overline{\mathrm{D}}} \mathrm{NPB}_{\overline{\mathrm{U}}} \\ & \mathrm{DT} \longrightarrow \mathrm{the} \; / \; \varepsilon \\ & \mathrm{JJ} \longrightarrow \mathrm{strong} \; / \; \varPsi \mathfrak{m} \\ & \mathrm{JJ} \longrightarrow \mathrm{north} \; / \; \mathfrak{k} \\ & \mathrm{NN} \longrightarrow \mathrm{wind} \; / \; \overline{\mathrm{M}} \end{split}$$

- 1-to-1 correspondence between nodes
- isomorphic derivation trees
- uniquely determined word alignment



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Bracketing Transduction Grammars

BTG (Wu, 1997):

- special form of SCFG
- only one nonterminal X
- nonterminal rules:

 $\begin{cases} X \to X_{[1]} \ X_{[2]} \ / \ X_{[1]} \ X_{[2]} & \text{monotone rule} \\ X \to X_{[1]} \ X_{[2]} \ / \ X_{[2]} \ X_{[1]} & \text{inversion rule} \end{cases}$

▶ preterminal rules where $e \in V_t \cup \{\epsilon\}$ and $f \in V_s \cup \{\epsilon\}$:

 $\Big\{X \to f \ / \ e \quad \text{lexical translation rules}$

SCFG Derivations

Generative Process:

- 1. Write down the **source** and **target** start symbols.
- 2. Choose a synchronous rule whose left-hand side is the left-most written down source nonterminal symbol.
- 3. Simultaneously rewrite the source symbol and its corresponding target symbol with the source and the target side of the rule, respectively.
- 4. Repeat step 2 while there are written down source and target nonterminal symbols.

$$\begin{cases} X \to X_{[1]} \ X_{[2]} \ / \ X_{[1]} \ X_{[2]} \\ X \to X_{[1]} \ X_{[2]} \ / \ X_{[2]} \ X_{[1]} \\ X \to k \ / \ k \end{cases} \quad k \in \{1, 2, 3\} \end{cases}$$

BTG Alignments



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Phrase-Based Models and SCFGs

SCFG Formalization of Phrase-Based Translation Models:

• \forall phrase pair $\langle \tilde{f}, \tilde{e} \rangle$, make rule

$$X \to \tilde{f} \ / \ \tilde{e}$$

make monotone rules

$$S
ightarrow S_{[1]} \; X_{[2]} \; / \; S_{[1]} \; X_{[2]}$$

 $S
ightarrow X_{[1]} \; / \; X_{[1]}$

make reordering rules

$$X \to X_{[1]} X_{[2]} / X_{[2]} X_{[1]}$$

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? Completeness ? Correctness ?

Hierarchical Phrase-Based Models

HPBM (Chiang, 2007):

- first tree-to-tree approach to perform better than phrase-based systems in large-scale evaluations
- discontinuous phrases
- Iong-range reordering rules
- formalized as synchronous context-free grammars

澳洲 是 与 北韩 有 邦交 的 少数 国家 之一 。 Aozhou shi yu Beihan you bangjiao de shaoshu guojia zhiyi . Australia is with North Korea have dipl. rels. that few countries one of .

Australia is one of the few countries that have diplomatic relations with North Korea.

HPBM: Motivations

Aozhou shi yu Beihan you bangjiao de shaoshu guojia zhiyi . Australia is with North Korea have dipl. rels. that few countries one of .

Typical Phrase-Based Chinese-English Translation:

[Aozhou] [shi]₁ [yu Beihan]₂ [you] [bangjiao] [de shaoshu guojia zhiyi] [.]
 [Australia] [has] [dipl. rels.] [with North Korea]₂ [is]₁ [one of the few countries] [.]

Chinese VPs follow PPs / English VPs precede PPs

yu X₁ you X₂ / have X₂ with X₁ ► Chinese NPs follow RCs / English NPs precede RCs

 X_1 de X_2 / the X_2 that X_1

translation of *zhiyi* construct in English word order

 X_1 zhiyi / one of X_1

HPBM: Example Rules

$S \rightarrow X_1 \ / \ X_1$	(1)
$S ightarrow S_1 X_2 / S_1 X_2$	(2)
$X ightarrow$ yu X_1 you X_2 / have X_2 with X_1	(3)
$X ightarrow X_1$ de X_2 / the X_2 that X_1	(4)
$X ightarrow X_1$ zhiyi / one of X_1	(5)
X → Aozhou / Australia	(6)
$X \rightarrow Beihan / N.$ Korea	(7)
$X \rightarrow she / is$	(8)
$X \rightarrow bangjiao / dipl.rels.$	(9)
X ightarrow shaoshu guojia / few countries	(10)

S

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 $S \rightarrow S_1 X_2 / S_1 X_2$





 $S \rightarrow S_1 X_2 / S_1 X_2$

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 $S \rightarrow X_1 / X_1$

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 $X \rightarrow Aozhou / Australia$

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 $X \rightarrow she / is$

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 $X \rightarrow X_1$ zhiyi / one of X_1

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 $X \rightarrow X_1$ de X_2 / the X_2 that X_1

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 $X \rightarrow yu X_1 you X_2 / have X_2 with X_1$

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 $X \rightarrow Beihan / N.$ Korea

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 $X \rightarrow bangjiao / dipl.rels.$

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 $X \rightarrow$ shaoshu guojia / few countries

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Summary

Synchronous Context-Free Grammars:

 formal model to synchronize source and target derivation processes

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- BTG alignment model
- HPB recursive reordering model

Additional topics (optional):

- Decoding SCFGs: Translation as Parsing
- Learning SCFGs from phrase tables

Synchronous Context-Free Grammars

SCFGs:

- CFGs in **two dimensions**
- synchronous derivation of isomorphic^a trees
- unbounded reordering preserving hierarchy

 $VB \rightarrow PRP_1 VB1_2 VB2_3 / PRP_1 VB2_3 VB1_2$ $VB2 \rightarrow VB_1 TO_2 / TO_2 VB_1 ga$ $TO \rightarrow TO_1 NN_2 / NN_2 TO_1$ $PRP \rightarrow he / kare ha$ $VB \rightarrow listening / daisuki desu$



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Weighted SCFGs

▶ rules $A \rightarrow \alpha / \beta$ associated with positive weights $\mathbf{w}_{A\rightarrow \alpha/\beta}$

• derivation trees $\pi = \langle \pi_1, \pi_2 \rangle$ weighted as

$$\mathbf{W}(\pi) = \prod_{A \to \boldsymbol{\alpha}/\beta \in G} \mathbf{w}_{A \to \boldsymbol{\alpha}/\beta}^{f_{A \to \boldsymbol{\alpha}/\beta}(\pi)}$$

probabilistic SCFGs if the following conditions hold

$$\mathbf{w}_{\mathcal{A}
ightarrow oldsymbol{lpha}/eta} \in [0,1] ext{ and } \sum_{oldsymbol{lpha},eta} \mathbf{W}_{\mathcal{A}
ightarrow oldsymbol{lpha}/eta} = 1$$

notice: SCFGs might well include rules of type

$$A \rightarrow \alpha/\beta_1 \dots A \rightarrow \alpha/\beta_k$$

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MAP Translation Problem

Maximum A Posterior Translation:

$$e^{\star} = \underset{e}{\operatorname{argmax}} p(e|f)$$

= $\underset{e}{\operatorname{argmax}} \sum_{\pi \in \Pi(f,e)} p(e,\pi|f)$
 $\Pi(f,e)$ is the set of synchronous derivation trees yielding $\langle f, e \rangle$

 Exact MAP decoding is NP-hard (Simaan, 1996; Satta and Peserico, 2005)

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Viterbi Approximation

Tractable Approximate Decoding:

$$e^{\star} = \operatorname{argmax}_{e} \sum_{\pi \in \Pi(f,e)} p(e,\pi|f)$$
$$\simeq \operatorname{argmax}_{e} \max_{\pi \in \Pi(f,e)} p(e,\pi|f)$$
$$= E(\operatorname{argmax}_{\pi \in \Pi(f)} p(\pi))$$

 $\Pi(f)$ is the set of synchronous derivations yielding f $E(\pi)$ is the target string resulting from the synchronous derivation π

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Translation as Parsing

$$\pi^{\star} = \operatorname*{argmax}_{\pi \in \Pi(\mathbf{f})} p(\pi)$$

Parsing Solution:

- compute the most probable derivation tree that generates *f* using the source dimension of the WSCFG
- 2. build the **translation string** *e* by applying the **target dimension** of the rules used in the most probable derivation
- most probable derivation computed in O(n³) using dynamic programming algorithms for parsing weighted CFGs
- transfer of algorithms and optimizations developed for CFG to SMT

Translation as Parsing: Illustration



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Weighted CFGs in Chomsky Normal Form

WCFGs:

- ▶ rules $A \rightarrow \alpha$ associated with positive weights $\mathbf{w}_{A\rightarrow \alpha}$
- derivation trees π weighted as

$$\mathsf{W}(\pi) = \prod_{A o \mathbf{\alpha} \in \mathcal{G}} \mathsf{w}_{A o \mathbf{\alpha}}^{f_{A o \mathbf{\alpha}}(\pi)}$$

probabilistic CFGs if the following conditions hold

$$\mathbf{w}_{\mathcal{A}
ightarrow oldsymbol{lpha}} \in [0,1] ext{ and } \sum_{oldsymbol{lpha}} \mathbf{w}_{\mathcal{A}
ightarrow oldsymbol{lpha}} = 1$$

WCFGs in CNF:

- \blacktriangleright rules in CFGs in Chomsky Normal Form: $\textbf{A} \rightarrow \textbf{BC}$ or $\textbf{A} \rightarrow \textbf{a}$
- equivalence between WCFGs and WCFGs in CNF
- no analogous equivalence holds for weighted SCFGs

Translation as Weighted CKY Parsing

Given a WSCFG G and a source string f:

1. project **G** into its source WCFG **G**

 $A \xrightarrow{w} \alpha \in G$ if $A \xrightarrow{w} \alpha/\beta \in G$ and $\forall A \xrightarrow{w'} \alpha/\beta' \in G \ w \ge w'$

- 2. transform G into its CNF G'
- 3. solve $\pi'^{\star} = \operatorname{argmax}_{\pi' \in \Pi_{G'}(f)} \mathbf{p}(\pi')$ with the CKY algorithm
- 4. revert π'^{\star} into π^{\star} , the derivation tree according to **G**
- 5. map π^* into its corresponding target tree π
- 6. read off the translation e from π

Weighted CKY Parsing

Dynamic Programming:

- recursive division of problems into subproblems
- optimal solutions compose optimal sub-solutions (Bellman's Principle)
- tabulation of subproblems and their solutions

CKY Parsing:

- ▶ subproblems: parsing substrings of the input string u₁...u_n
- solutions to subproblems tabulated using a chart
- bottom up algorithm starting with derivation of terminals
- $O(n^3|G|)$ time complexity
- widely used to perform statistical inference over random trees

Weighted CKY Parsing

Problem:
$$\pi^* = \operatorname{argmax}_{\pi \in \Pi_G(u=u_1...u_n)} p(\pi)$$

DP chart:

 $M_{i,k,A}$ = maximum probability of $A \Rightarrow^* u_{i+1,k}$

base case, k - i = 1:

$$\forall \ 1 \leq i \leq n \ M_{i-1,i,A} = \mathbf{w}_{A \to u_i})$$

• inductive case, k - i > 1:

$$M_{i,k,A} = \max_{B,C,i < j < k} \{ \mathbf{w}_{A \to B \ C} \times M_{i,j,B} \times M_{j,k,C} \}$$

▶ best derivation built by storing *B*, *C*, and *j* for each $M_{i,k,A}$

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Weighted CKY Parsing



Weighted CKY Pseudo-Code

1: $\forall A, 0 \leq i, j \leq n \ M_{i,i,A} = 0;$ 2: for i = 1 to *n* do {base case: substrings of length 1} 3: $M_{i-1,i,A} = \mathbf{w}_{A \to u_i};$ {inductive case: substrings of length > 1} 4: for l = 2 to *n* do {*l*: length of the subtring} **for** i = 0 to n - l **do** {*i*: start position of the substring} 5: k = i + l; {k: end position of the substring} 6: **for** j = i + 1 to k - 1 **do** {j: split position} 7: for $\forall A \rightarrow BC$ do 8: $q = \mathbf{w}_{A \to BC} \times M_{i,i,B} \times M_{i,k,C}$ 9: if $q > M_{i,k,A}$ then 10: $M_{i,k,A} = q;$ 11: {backpointers to build derivation tree} 12: $D_{i,k,A} = \langle i, B, C \rangle;$

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Parsing SCFG and Language Modelling

Viterbi Decoding of WSCFGs:

- focus on most probable derivation of source (ignoring different target sides associated with the same source side)
- derivation weights do not include language models scores
- ? HOW TO EFFICIENTLY COMPUTE TARGET LANGUAGE MODEL SCORES FOR POSSIBLE DERIVATIONS ?

Approaches:

- 1. **rescoring**: generate *k*-best candidate translations and rerank *k*-best list with LM
- 2. **online**: integrate target *m*-gram LM scores into dynamic programming parsing
- 3. **cube pruning** (Huang and Chiang, 2007): rescore *k*-best sub-translations at each node of the parse forest

Online Translation

 Bàowēier yǔ
 Shālóng jǔxíng le huitán

 Powell
 with Sharon hold [past] meeting

 "Powell held a meeting with Sharon"

\mathbf{S}	\rightarrow	$NP^{(1)} VP^{(2)}$,	$NP^{(1)} VP^{(2)}$
VP	\rightarrow	$PP^{(1)} VP^{(2)}$,	$VP^{(2)} PP^{(1)}$
NP	\rightarrow	$B\dot{a}owar{e}ier,$	Powell
VP	\rightarrow	jŭxíng le huìtán,	held a meeting
\mathbf{PP}	\rightarrow	yŭ Shālóng,	with Sharon

Online Translation: parsing of the source string and building of the corresponding subtranslations **in parallel**

$$\frac{PP_{1,3}:(w_1,t_1) \ VP_{3,6}:(w_2,t_2)}{VP_{1,6}:(w \times w_1 \times w_2,t_2t_1)}$$

- w₁, w₂: weights of the two antecedents
- w: weight of the synchronous rule
- ▶ t₁, t₂: translations



LM Online Integration (Wu, 1996)

Bigram Online Integration:

 $\frac{PP_{1,3}^{with*Sharon}:(w_1,t_1) \ VP_{3,6}^{held*talk}:(w_2,t_2)}{VP_{1,6}^{held*Sharon}:(w \times w_1 \times w_2 \times p_{LM}(with|talk),t_2t_1)}$ Bush held a talk with Sharon
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Cube Pruning (Huang and Chiang, 2007)



Beam Search:

- at each step in the derivation, keep at most k items integrating target subtranslations in a beam
- enumerate all possible combinations of LM items
- extract the k-best combinations

Cube Pruning:

- get k-best LM items without without computing all possible combinations
- approximate beam search: prone to search errors (in practice, much less significant than efficient decoding)

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Cube Pruning

Heuristic Assumption:

- best adjacent items lie towards the upper-left corner
- part of the grid can be pruned without computing its cells





k-best parsing (Huang and Chiang, 2005)

- a priority queue of candidates
- extract the best candidate



$({\rm VP} \; {}^{\rm held}_{3,6} \star {}^{\rm meeting})$
$(\mathrm{VP} \; _{3,6}^{\mathrm{held} \; \star \; \mathrm{talk}})$
$(VP_{3,6}^{hold \star conference})$

	1.0	3.0	8.0
1.0	2.5	9.0	9.5
1.1	2.4	9.5	9.4
3.5	5.1	17.0	12.1

k-best parsing (Huang and Chiang, 2005)

- a priority queue of candidates
- extract the best candidate
- push the two successors



$({\rm VP}~^{\rm held~\star~meeting}_{3,6})$
$(\mathrm{VP}_{3,6}^{\mathrm{held}\;\star\;\mathrm{talk}})$
$\mathrm{VP}_{3,6}^{\mathrm{hold} \star \mathrm{conference}})$

	1.0	3.0	8.0
1.0	2.5	9.0	9.5
1.1	2.4	9.5	9.4
3.5	5.1	17.0	12.1

k-best parsing (Huang and Chiang, 2005)

- a priority queue of candidates
- extract the best candidate
- push the two successors



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$({\rm VP} \; {}^{\rm held}_{3,6} \star {}^{\rm meeting})$	Ι.
$(\mathrm{VP}_{3,6}^{\mathrm{held}\;\star\;\mathrm{talk}})$].
${ m VP}^{ m hold}_{ m 3,6} \star { m conference})$	3.

	1.0	3.0	8.0
1.0	2.5	9.0	9.5
1.1	2.4	9.5	9.4
3.5	5.1	17.0	12.1

Cube Pruning: Pseudo-Code

To efficiently compute a small corner of the grid:

- push cost of grid cell 1, 1 onto priority queue
- repeat j
 - 1. extract best cell from queue
 - 2. push costs of best cell's neighbours onto queue
- until k cells have been extracted (other termination conditions are possible)

```
1: function CUBE(F)
                                                   \triangleright the input is a forest F
 2:
          for v \in F in (bottom-up) topological order do
 3:
               KBEST(v)
          return D_1(TOP)
 4:
     procedure KBEST(v)
 5:
 6:
          cand \leftarrow \{ \langle e, \mathbf{1} \rangle \mid e \in IN(v) \} \triangleright \text{ for each incoming } e
 7:
          HEAPIFY(cand)
                                         > a priority queue of candidates
          buf \leftarrow \emptyset
 8:
 9:
          while |cand| > 0 and |buf| < k do
10:
               item \leftarrow POP-MIN(cand)
11:
               append item to buf
12:
               PUSHSUCC(item, cand)
13:
          sort buf to \mathbf{D}(v)
14: procedure PUSHSUCC(\langle e, j \rangle, cand)
15:
          e \text{ is } v \rightarrow u_1 \dots u_{|e|}
16:
          for i in 1 \dots |e| do
17:
               \mathbf{j}' \leftarrow \mathbf{j} + \mathbf{b}^i
18:
               if |\mathbf{D}(u_i)| > j'_i then
                    PUSH(\langle e, j' \rangle, cand)
19:
```

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Summary

Translation As Parsing:

- Viterbi Approximation
- Weighted CKY Parsing
- Online LM Integration and Cube Pruning

Next Session:

Learning SCFGs and Hiero

Hierarchical Phrase-Based Models

Hiero (Chiang, 2005, 2007):

 SCFG of rank 2 with only two nonterminal symbols

- discontinuous phrases
- long-range reordering rules



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Hiero Synchronous Rules



Rule Extraction:

- a word-aligned sentence pair
- b extract initial phrase pairs
- c replace sub-phrases in phrases with symbol *X*

Glue Rules:

$$S \rightarrow S_1 X_2 / S_1 X_2 \ S \rightarrow X_1 / X_1$$

Rule Filtering:

- limited length of initial phrases
- no adjacent nonterminals on source
- at least one pair of aligned words in non-glue rules

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 $X \rightarrow 5$ 北 韩 有 邦交, have diplomatic relations with North Korea

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 $X \rightarrow 5$ 北 韩 有 邦交, have diplomatic relations with North Korea

X → 邦交, diplomatic relations

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 $X \rightarrow 5$ 北 韩 有 邦交, have diplomatic relations with North Korea

X → 邦交, diplomatic relations X → 北 韩, North Korea



 $X \rightarrow 5$ 北 韩 有 邦交, have diplomatic relations with North Korea

X → 邦交, diplomatic relations X → 北 韩, North Korea

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 $X \rightarrow 5$ 北 韩 有 邦交, have diplomatic relations with North Korea

X → 邦交, diplomatic relations X → 北 韩, North Korea X → 与 X1 有 X2,

have X_2 with X_1



 $X \rightarrow 5$ 北 韩 有 邦交, have diplomatic relations with North Korea

X → 邦交, diplomatic relations X → 北 韩, North Korea X → 与 X1 有 X2,

have X_2 with X_1

Hiero: Log-Linear Parametrization

Scoring Rules:

$$\mathcal{S}(\mathsf{A} \to \gamma) = \lambda \cdot \mathsf{h}$$

• $h(A \rightarrow \gamma)$: feature representation vector $\in \mathbb{R}^m$

- λ : weight vector $\in \mathbb{R}^m$
- $h_r(A \rightarrow \gamma)$: value of the *r*-th feature
- λ_r: weight of the r-th feature

Scoring Derivations:

$$\mathcal{S}(\pi) = \lambda_{LM} \log p(E(\pi)) + \sum_{\langle Z \to \gamma, i, j \rangle \in \pi} \mathcal{S}(Z \to \gamma)$$

- derivation scores decompose into sum of rule scores
- $p(E(\pi))$ is the LM score computed while parsing

Hiero: Feature Representation

Word Translation Features:

$$h_1(X o {oldsymbol lpha}/{oldsymbol eta}) = \log p(\mathcal{T}_{oldsymbol eta} | \mathcal{T}_{oldsymbol lpha})$$

 $h_2(X \to \alpha/\beta) = \log p(\mathcal{T}_{\alpha}|\mathcal{T}_{\beta})$

Word Penalty Feature:

$$h_3(X \to \alpha/\beta) = -|\mathcal{T}_\beta|$$

Synchronous Features:

$$h_4(X \to \alpha/\beta) = \log p(\beta|\alpha)$$

$$h_5(X \to \frac{\alpha}{\beta}) = \log p(\frac{\alpha}{\beta})$$

Glue Penalty Feature:

$$h_6(S \rightarrow \frac{S_1X_1}{S_1X_1}) = -1$$

Phrase Penalty Feature:

$$h_7(X \to \alpha/\beta) = -1$$

λ_i tuned on dev set using MERT

Summary

Hiero:

- Rule Extraction
- Log-Linear Parametrization

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Feature Representation